(Auto) Verus Building Software that You Can Trust

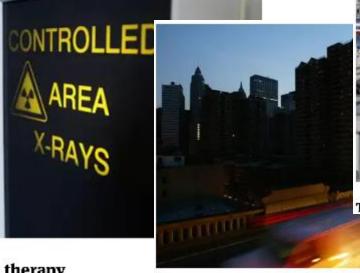
Chenyuan Yang



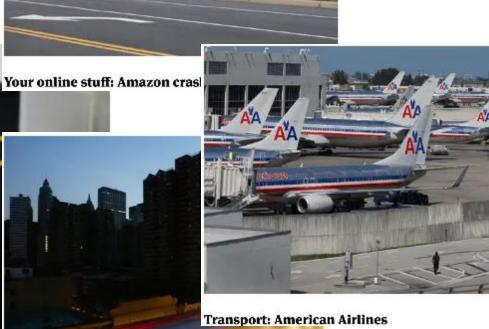
Software correctness is critical



Space: Ariane 5

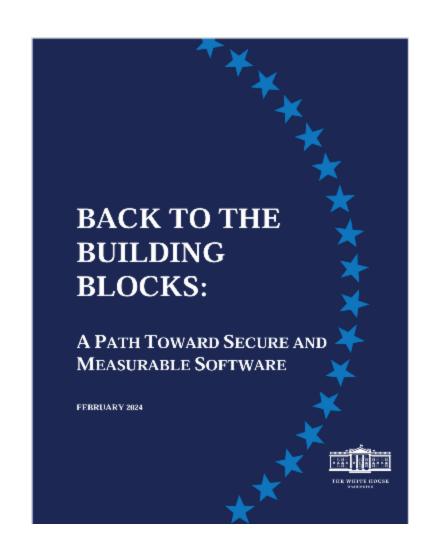


Medical: Radiation therapy



Infrastructure: The north-east US power outage

For reliability & security, developers ...

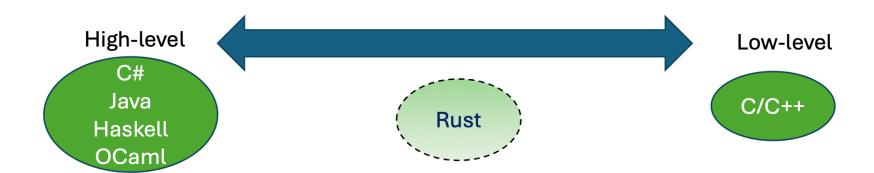


Use memory-safe programming languages, such as Rust ...

Use formal verification for the core components...

Why Rust?

- Ownership: Every value has a unique "owner."
- Borrowing: You can borrow a value, but there are strict rules.
 - One mutable borrow OR multiple immutable borrows.
- Lifetimes: The compiler ensures that references don't outlive the data they point to.



Rust borrow example

```
fn main() {
  let s1 = String::from("hello");
  // s1's ownership is MOVED to the function
  takes_ownership(s1)
  // This line would cause a compiler error!
  // println!("s1 is {}", s1);
  let s2 = String::from("world");
  // s2 is BORROWED by the function
  borrows_immutably(&s2);
  // s2 is still valid here because it was borrowed
  println!("s2 is still {}", s2);
```

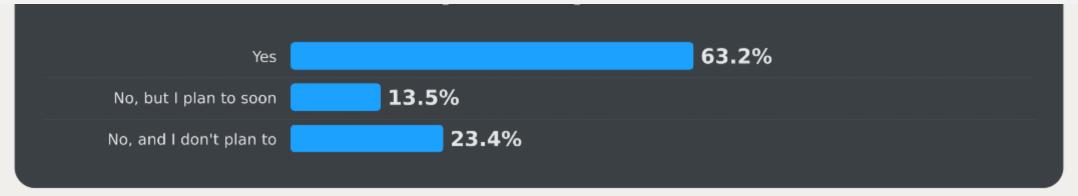
```
fn takes_ownership(some_string: String) {
    println!("{}", some_string);
} // `some_string` is dropped here

fn borrows_immutably(some_string: &String) {
    println!("{}", some_string);
} // `some_string` is not dropped
```

For convenience, developers ...

Using AI assistants or coding agents is already the trend!

Do you currently use AI tools in your development process?

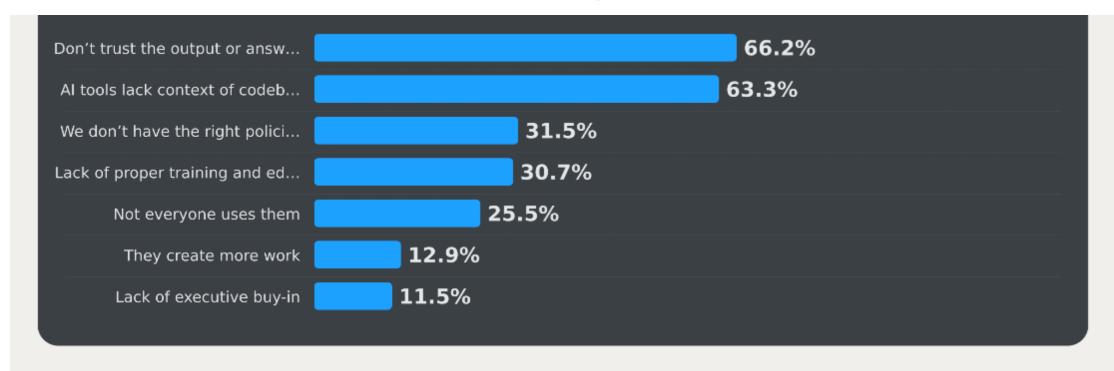




Source: survey.stackoverflow.co/2024
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AI? Reliability & Security?

66.2% Don't trust the output or answers of Al





Source: survey.stackoverflow.co/2024 Data licensed under Open Database License (ODbL)

How can we trust Al-generated code?

- Software testing
 - To expose bugs in code
 - Active research in AI for testing and testing for AI

But it cannot make sure that there is no bug (



- Software verification
 - To mathematically prove important properties of code
 - To prove that there is no bug!

Why not formally verify software?

Can I verify the software itself instead of a model of it?



Can I not learn a new language to write spec/proof?



How long does it take the verifier to run?



How fast is the verified software?



You should try Verus!

Can I verify the software itself instead of a model of it?



Can I not learn a new language to write spec/proof?



How long does it take the verifier to run?



How fast is the verified software?



"Verus is a tool for verifying the correctness of code written in Rust. Developers write specifications of what their code should do, and Verus statically checks that the executable Rust code will always satisfy the specifications for all possible executions of the code"



GitHub
https://github.com/verus-lang/verus

verus

Verified Rust for low-level systems code

Primary language: Rust . License: MIT license

Verus is already used for various systems projects

- Persistent-memory log, key-value store for Azure Storage
- VeriSMo security module
 OSDI Best Paper
- Concurrent memory allocator (CMU)
- Atmosphere microkernel (U-Utah)
- Anvil Cluster Management (*U-Illinois*, *U-Wisconsin*, *VMware*)

OSDI Best Paper

"Verus: A Practical Foundation for Systems Verification"

SOSP Best Artifact

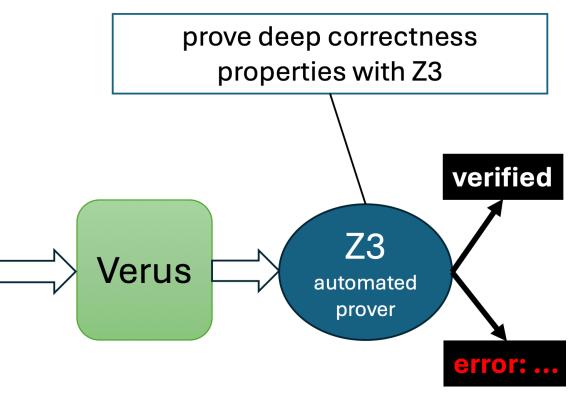
"Linear Types for Large-Scale Systems Verification"

OOPSLA Best Paper

How Verus works?

express specs, proofs in Rust

```
fn binary_search(v: &Vec<x64>, key: u64)
                               -> (ret: usize)
   requires
     forall|i:int, j:int / ... ==> v[i] <= v[j],
   ensures
     key == v[ret], ../.
   let mut left = 0;
   let mut right = \sqrt{.len()} - 1;
   while left != ri/ght
     invariant
       right < v.len(), ...
```



How Verus works?

```
fn binary_search(v: &Vec<u64>, key: u64) -> (ret: usize)
    let mut left: usize = 0;
    let mut right: usize = v.len() - 1;
    while left != right
        let mid = left + (right - left) / 2;
        if v[mid] < key {</pre>
            left = mid + 1;
        } else {
            right = mid;
```

How should we describe the functionality of binary_search?

Pre-Condition requires Post-Condition ensures

"The input **v** is a sorted vector and it has the value **key** we want to find"

"The output **ret** is a valid index **i** such that **v[i]** equals the **key**"

But ...

Can Al generate Verus specifications and proofs for me?

```
ţζ Ⅲ ···
misc > src > ® main.rs > rust-analyzer > ♥ binary_search
 16
 17
          fn binary_search(v: &Vec<u64>, k: u64) -> (r: usize)
 18
 19
               forall |i: int, j: int | 0 <= i < j < v@.len() ==> v[i] <= v[j],
               exists |i: int| 0 <= i < v@.len() && v[i] == k,
 20
 21
 22
               0 <= r < v@.len(),</pre>
               v[r as int] == k,
 23
 24
 25
               let mut i1: usize = 0;
 26
               let mut \overline{i2}: usize = v.len() - 1;
 27
 28
               while \underline{i1} != \underline{i2}
 29
 30
                  let ix: usize: usize = i1 + (i2 - i1) / 2;
 31
                  if v[ix] < k {
 32
                      i1 = ix + 1;
 33
                   } else {
 34
                       i2 = ix;
 35
 36
 37
 38
                                                                                                                                             ∨ ≣ A ⊕ … ^ ×
PROBLEMS OUTPUT DEBUG CONSOLE TERMINAL PORTS COMMENTS
                                                                                                                              Verus Copilot
```

Can GPT-4 prove binary-search?

```
fn binary search(v: &Vec<u64>, key: u64) -> (ret: usize)
requires
    forall | i: int, j: int | 0 <= i < j < v@.len() ==> v[i] <= v[j],
    exists |i:int| 0 <= i < v@.len() && key == v[i],
ensures
    ret < v.len(), key == v[ret as int],</pre>
    let mut left: usize = 0;
    let mut right: usize = v.len() - 1;
    while left != right
        let mid = left + (right - left) / 2;
        if v[mid] < key {</pre>
            left = mid + 1;
        } else {
            right = mid;
    left
```

"You are a Verus expert.

Please add proof annotations for the following code, so that Verus can prove the function Implementation satisfies the function specification."







??

Lack of knowledge: syntax

```
fn binary_search(v: &Vec<u64>, key: u64) -> (ret: usize)
requires
    forall | i: int, j: int | 0 <= i < j < v@.len() ==> v[i] <= v[j],
    exists | i: int | 0 <= i < v@.len() && key == v[i],
ensures
    ret < v.len(), key == v[ret as int],</pre>
    let mut left: usize = 0;
    let mut right: usize = v.len() - 1;
    while left != right
      invariant
        exists |i: usize | 0 <= i < v@.len() && key == v[i],
    { ... }
```

Verus error: mismatched type, expecting `int' yet getting `usize'

Lack of skills: loop invariants

```
fn binary search(v: &Vec<u64>, key: u64) -> (ret: usize)
requires
    forall | i: int, j: int | 0 <= i < j < v@.len() ==> v[i] <= v[j],
    exists |i:int| 0 \le i < v@.len() && key == v[i],
ensures
    ret < v.len(), key == v[ret as int],</pre>
                                                 Missing invariant:
    let mut left: usize = 0;
    let mut right: usize = v.len() - 1;
                                                  right < v@.len()
    while left != right
      invariant
        forall |i:int, j:int| 0 <= i < j < v@.len() ==> v[i] <= v[j],
        exists | i: int | left <= i <= right && key == v[i],
        let mid = left + (right - left) / 2;
        if v[mid] < key {</pre>
            left = mid + 1;
        } else {
            right = mid;
```

Lack of skills: loop invariants

```
fn binary search(v: &Vec<u64>, key: u64) -> (ret: usize)
requires
    forall | i: int, j: int | 0 <= i < j < v@.len() ==> v[i] <= v[j],
    exists |i:int| 0 \le i < v@.len() && key == v[i],
ensures
    ret < v.len(), key == v[ret as int],</pre>
                                                  Missing invariant:
    let mut left: usize = 0;
    let mut right: usize = v.len() - 1
                                                   right < v@.len()
    while left != right
      invariant
        forall |i:int, j:int| 0 <= i < j < v@.len() ==> v[i] <= v[j],
        exists | i: int | left <= i <= right && key == v[i],
        let mid = left + (right - left) /
                                              Verus error:
        if v[mid] < key {</pre>
                                              1. Function postconditions not satisfied
            left = mid + 1;
                                              2. invariant not satisfied at the end of loop
        } else {
                                              3. Precondition, mid < v@.len(), of v[mid] not satisfied
            right = mid;
```

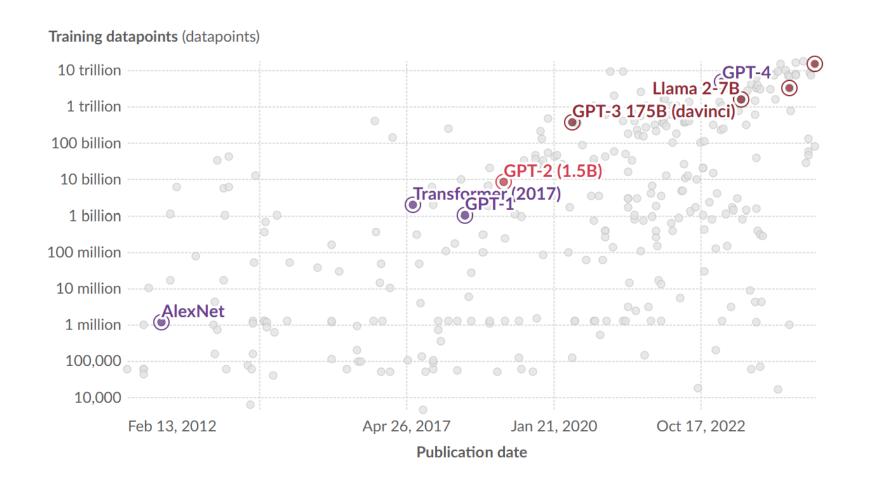
Lack of strategy: debugging, prioritization, ...

```
fn binary search(v: &Vec<u64>, key: u64) -> (ret: usize)
requires
    forall | i: int, j: int | 0 <= i < j < v@.len() ==> v[i] <= v[j],
    exists |i:int| 0 \le i < v@.len() && key == v[i],
ensures
    ret < v.len(), key == v[ret as int],</pre>
    let mut left: usize = 0;
    let mut right: usize = v.len() - 1;
    while left != right
      invariant
        forall |i:int, j:int| 0 <= i < j < v@.len() ==> v[i] <= v[j],
        exists |i: int| left <= i <= right && key == v[i],
        let mid = left + (right - left) /
                                             Verus error:
        if v[mid] < key {</pre>
                                              1. Function postconditions not satisfied
            left = mid + 1;
                                              2. invariant not satisfied at the end of loop
        } else {
                                              3. Precondition, mid < v@.len(), of v[mid] not satisfied
            right = mid;
```

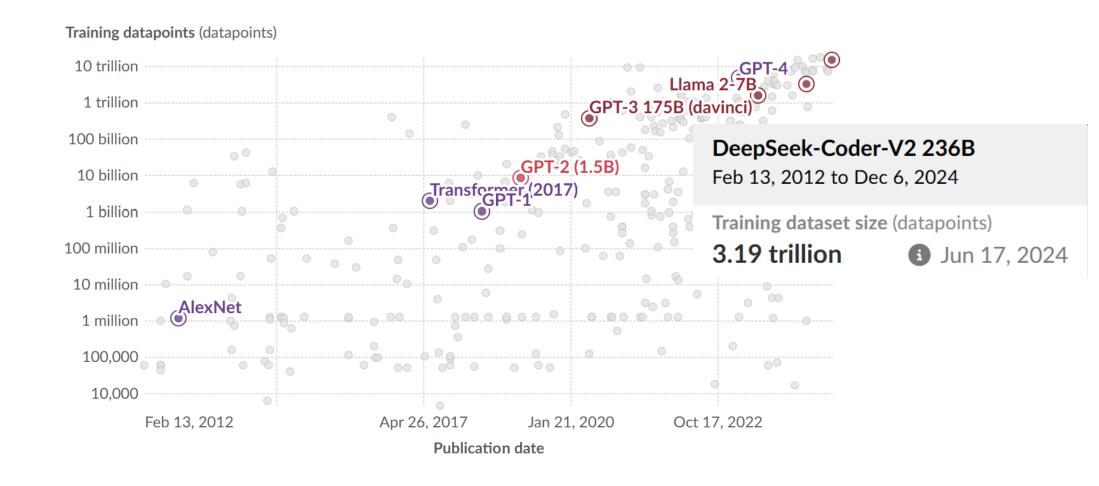
How to teach AI proof knowledge, skills, strategies?

Not enough data!

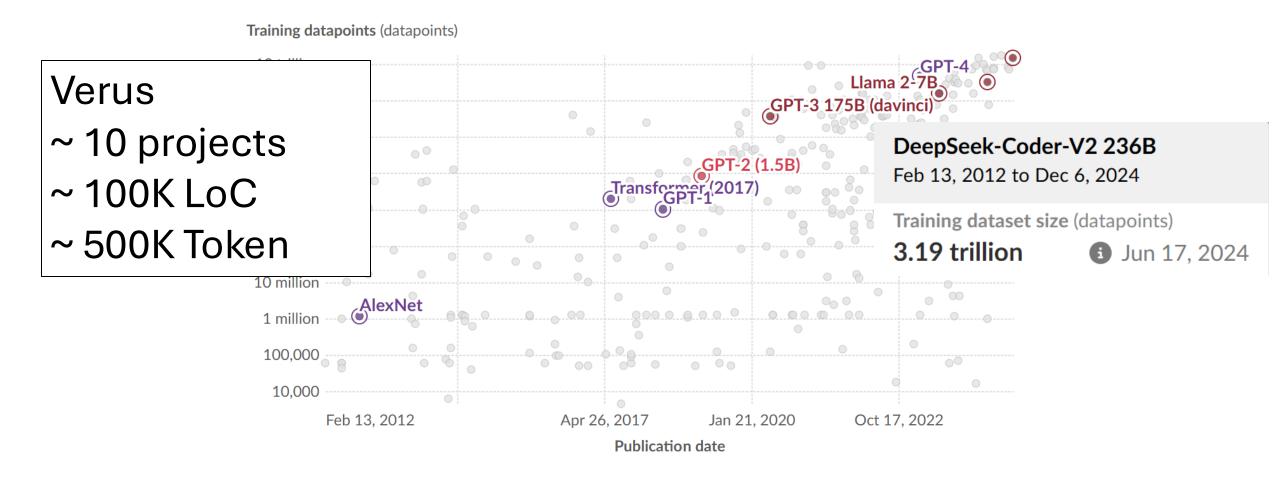
The amount of data used to train models



The amount of data used to train models



The amount of Verus data available



AutoVerus:

Automated Proof Generation for Rust Code

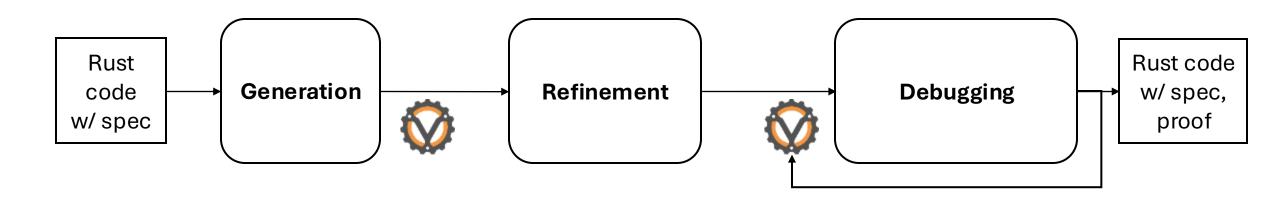
An agent framework supporting LLM through prompts, workflow, compiler & formal methods

OOPSLA 2025

How to teach LLMs to write Verus proof?

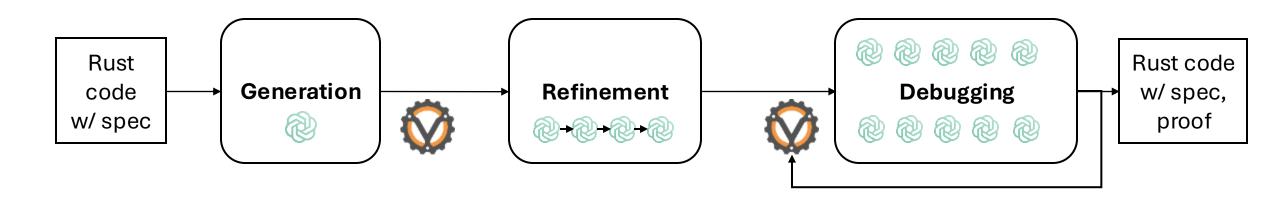
A workflow that mimics human experts' methodology

"The proof development of human experts is an iterative process of repeatedly running Verus, checking and prioritizing Verus errors, developing and editing proof to fix them." – Interview of multiple co-authors of the Verus paper



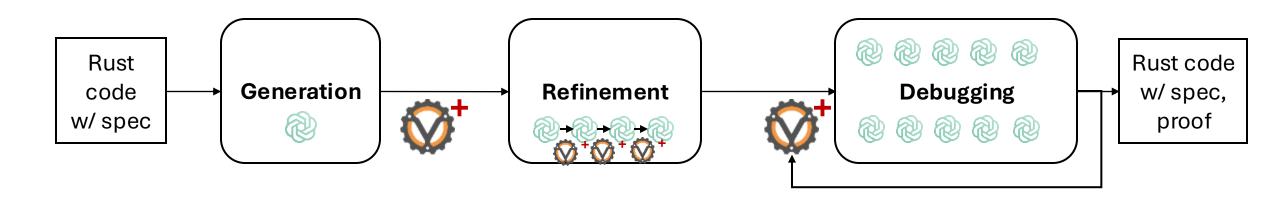
How to teach LLMs to write Verus proof?

- A workflow that mimics human experts' methodology
- An extensible network of GPT agents w/ Verus knowledge, skills

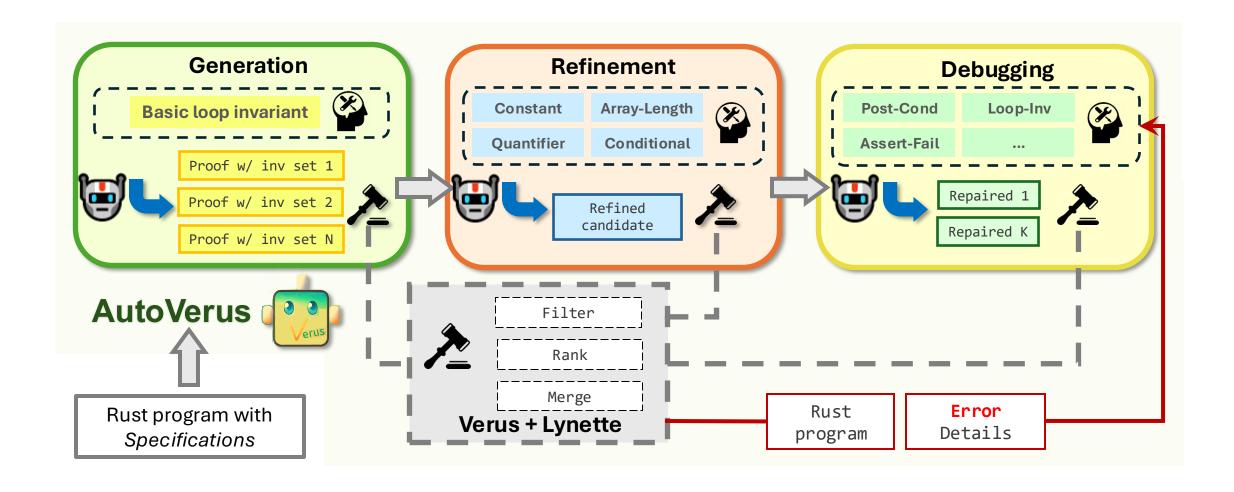


How to teach LLMs to write Verus proof?

- A workflow that mimics human experts' methodology
- An extensible network of GPT agents w/ Verus knowledge, skills
- Proof discipline, ranking, merging, ... supported by formal methods



AutoVerus



Lack of Discipline

LLM may cheat on the verification

error: invariant not satisfied after loop



`assume` makes the verification pass, but by adding assumption

Lack of Discipline

• LLM may change the executable code

```
while (i < v.len())
{
   if (v[i] <= e) {
     result.push(v[i]);
   }
   i = i + 1;
}</pre>
```



```
while (i < v.len())
{
    if (v[i] > e) {
       result.push(v[i]);
    }
    i = i + 1;
}
```

We need discipline!

We DO want to unleash the creativity of LLM

But

- We do NOT want LLM to make arbitrary changes to the input
- We need an effective way to search among many creative outputs

How to add discipline to LLM?

Using Verus + Lynette to

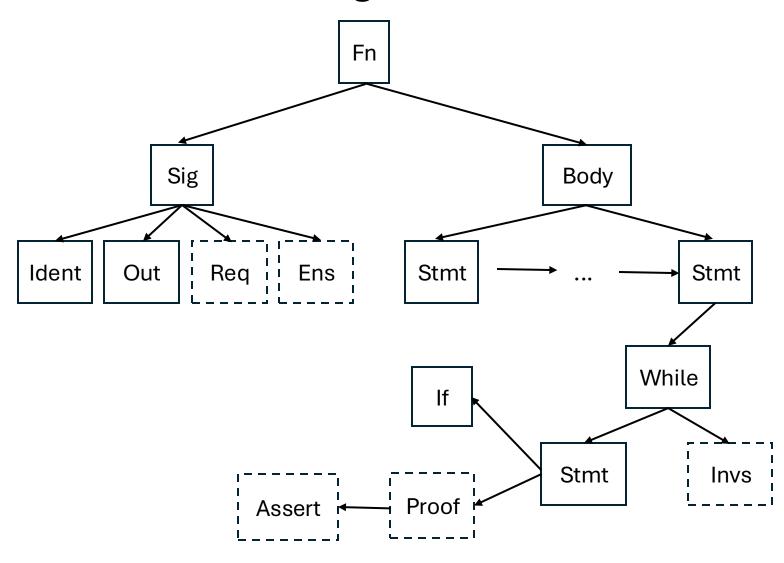
- Rank
- Merge
- Filter

all LLM outputs

The exact ranking/merging/filtering policy is skipped from this talk

Discipline - Lynette, the Verus Source Forger

```
fn remove all greater(v: Vec<i32>, i: T)
  -> (result: Vec<i32>)
  requires
    forall |k1:int,k2:int|
      0 <= k1 < k2 < v.len() ==> v[k1] != v[k2]
  ensures
    forall |k:int| 0 <= k < result.len()</pre>
==>
      result[k] <=
e&&v@.contains(result[k]),
    forall |k:int| 0 <= k < v.len()
      &&
v[k]<=e==>result@.contains(v[k]),
  let mut i: usize = 0;
  let vlen = v.len();
  let mut result: Vec<i32> = vec![];
  while (i < v.len())</pre>
    if (v[i] <= e) {</pre>
      result.push(v[i]);
    i = i + 1;
    proof {
      assert(...);
      . . .
  result
```



Discipline - Lynette, the Verus Source Forger

- Detecting "unsafe" changes
 - AST-level comparison
 - Bottom-line: Generate same executable code
 - By erasing all ghost code
 - Then comparing the rest of the code
 - Conditional
 - Spec function
 - Pre/post condition
 - Assumption

Expertise – Error-Fix Action Table

Error Type					
Function postcondition not satisfied					
Function precondition not satisfied					
Function precondition not satisfied (Vector Length Violation)					
Loop invariant not satisfied at end of loop body					
Loop invariant not satisfied before the loop body					
Assertion failed					
Arithmetic overflow/underflow					
Type error					
Misc. verus syntax error					
···					

Expertise – Error-Fix Action Table

Error Type	Fix Actions				
Postcondition not	Add the proof blocks related to the post-condition at the exit point				
satisfied	Modify the existing loop invariants				
Precondition not satisfied	Add the assertions related to the pre-condition just before the invocation of the function				
Precondition not satisfied - Vector Length	add loop invariants/asserts for the array: 1. an invariant that specify the array length (i.e., A.len() ==); 2. an invariant about the array index not under bound (e.g., k >= 0)				
Invariant not satisfied at end of loop body	the end of the loop				
Invariant not satisfied before loop	One LLM agent for each fix action nt before the loop body e it correct				
	Add the failed invariant to all the loops before the failed loop				
	Delete the failed loop invariant				
	Add the necessary assertions before the failed assertion				
Assertion failed	Add appropriate loop invariants to ensure the assertion holds true				
	Fix the assertion error for the following code by using existing lemma functions				
	Fix the assertion error for the following code by creating the helper proof functions				

Repair: Post-Condition Not Satisfied

Your mission is to **fix the post-condition not satisfied error** for the following code. Basically, you should add the proof blocks related to the post-condition at the exit point, or modify the existing loop invariants to make them work for the post-condition

There are two general fixes for the "post-condition not satisfied" error

```
pub fn filter(x: &Vec<u64>, y: &mut Vec<u64>)
requires
    old(y).len() == 0,
ensures
    y@ == x@.filter(|k:u64| k%3 == 0),
    let mut i: usize = 0;
   let xlen = x.len();
   while (i < xlen)
   invariant
        i <= xlen.
        y@ == x@.take(i as int).filter(|k:u64| k%3 ==
0),
        if (x[i] % 3 == 0) {
            y.push(x[i]);
        i = i + 1;
```

```
pub fn filter(x: &Vec<u64>, y: &mut Vec<u64>)
requires
    old(y).len() == 0,
ensures
    y@ == x@.filter(|k:u64| k%3 == 0),
    let mut i: usize = 0;
   let xlen = x.len();
   while (i < xlen)
   invariant
        i <= xlen.
        y@ == x@.take(i as int).filter(|k:u64| k%3 ==
0),
        if (x[i] \% 3 == 0) {
            y.push(x[i]);
        i = i + 1:
    proof {
        assert(y@ == x@.filter(|k:u64| k%3 == 0));
    } // Added by AI
```

Benchmark Construction: Verus-Bench

- No existing Verus proof generation benchmark
- We translated three verification-related benchmark in other languages (C, Dafny) into Verus
 - CloverBench^[1], Diffy^[2], MBPP^[3]
 - Misc is collected from Verus tutorials
- The first benchmark designed for Verus proof generation

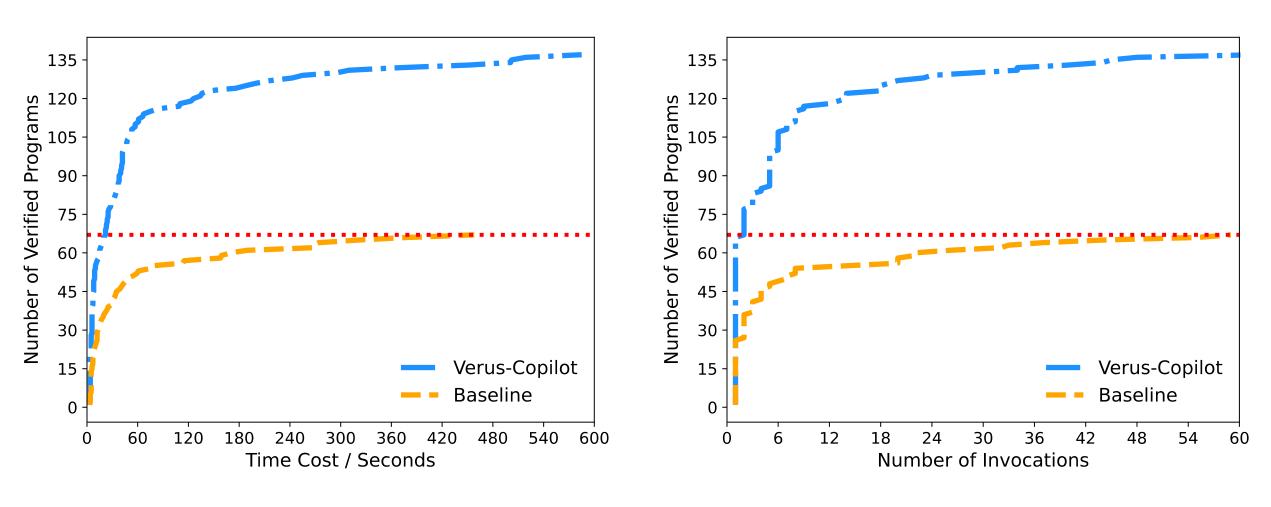
Benchmark Sources	CloverBench	Diffy	MBPP	Misc	Total
# of Proof Tasks	11	38	78	23	150
Executable LOC Specification LOC	175 80	951 265	1,333 700	390 207	2,849 1,252

^[1] Sun, Chuyue, et al. "Clover: Clo sed-Loop Ver ifiable Code Generation." International Symposium on Al Verification. Cham: Springer Nature Switzerland, 2024.

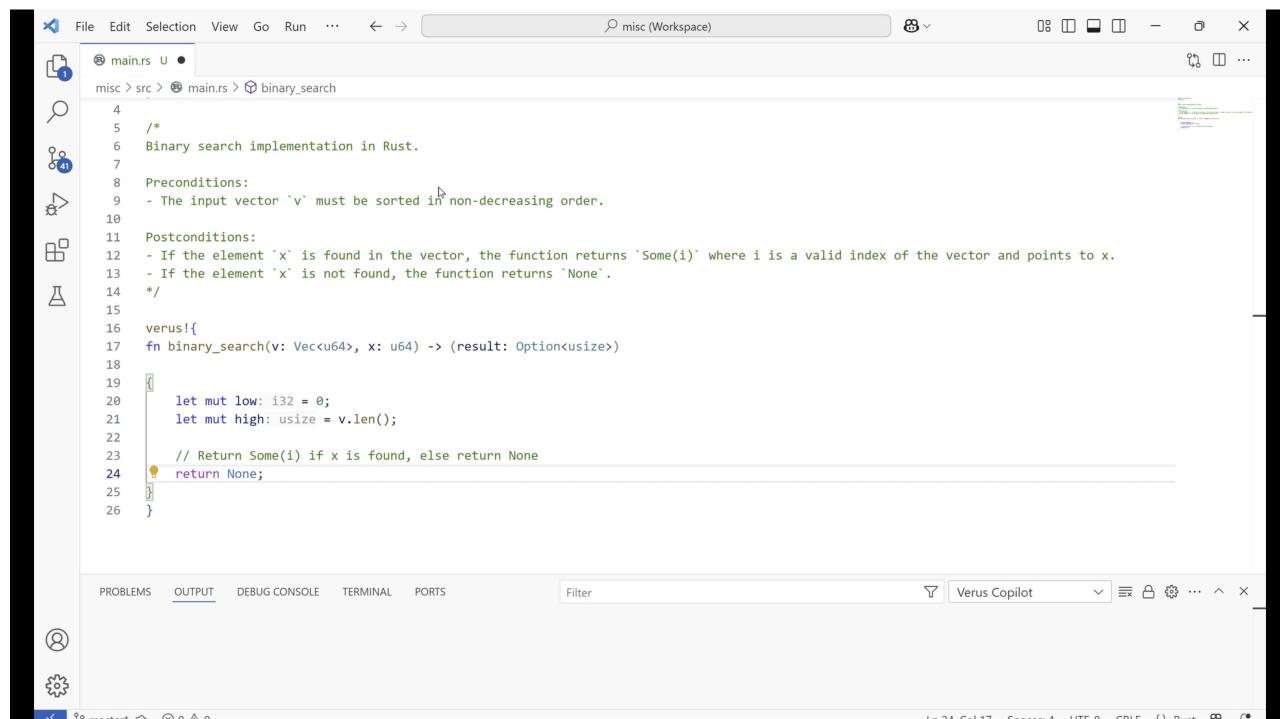
^[2] Chakraborty, Supratik, Ashutosh Gupta, and Divyesh Unadkat. "Diffy: Inductive reasoning of array programs using difference invariants.", CAV 2021,

^[3] Misu, Md Rakib Hossain, et al. "Towards ai-assisted synthesis of verified dafny methods." Proceedings of the ACM on Software Engineering 1.FSE (2024)

Results Based on Time and Invocation



Much better performance than directly invoking LLMs



Recap

- Verus could make sure that AI-generated code is 100% correct!
 - Precondition
 - Postcondition
 - Proof annotation like loop invariants, assertions, etc
- All could also help you to complete the proof!
 - AutoVerus: https://github.com/microsoft/verus-proof-synthesis
 - Verus-Copilot: https://github.com/microsoft/verus-copilot-vscode
 - VSCode extension